

Coarray Status

Dan Nagle

PGAS

- Partitioned Global Address Space
- Examples include:
 - UPC (upc.gwu.edu)
 - Titanium (titanium.cs.berkeley.edu)
- Original Paper at <ftp://ftp.numerical.rl.ac.uk/pub/reports/nrRAL98060.ps.gz>

Other Sources

- WG5 NI824.pdf (John Reid's summary)
(www.nag.co.uk/sc22wg5) (see also NI830)
- Rice University (caf.rice.edu)
- FDIS J3/10-007r1.pdf (www.j3-fortran.org)
- Modern Fortran Explained by Metcalf, Reid, Cohen

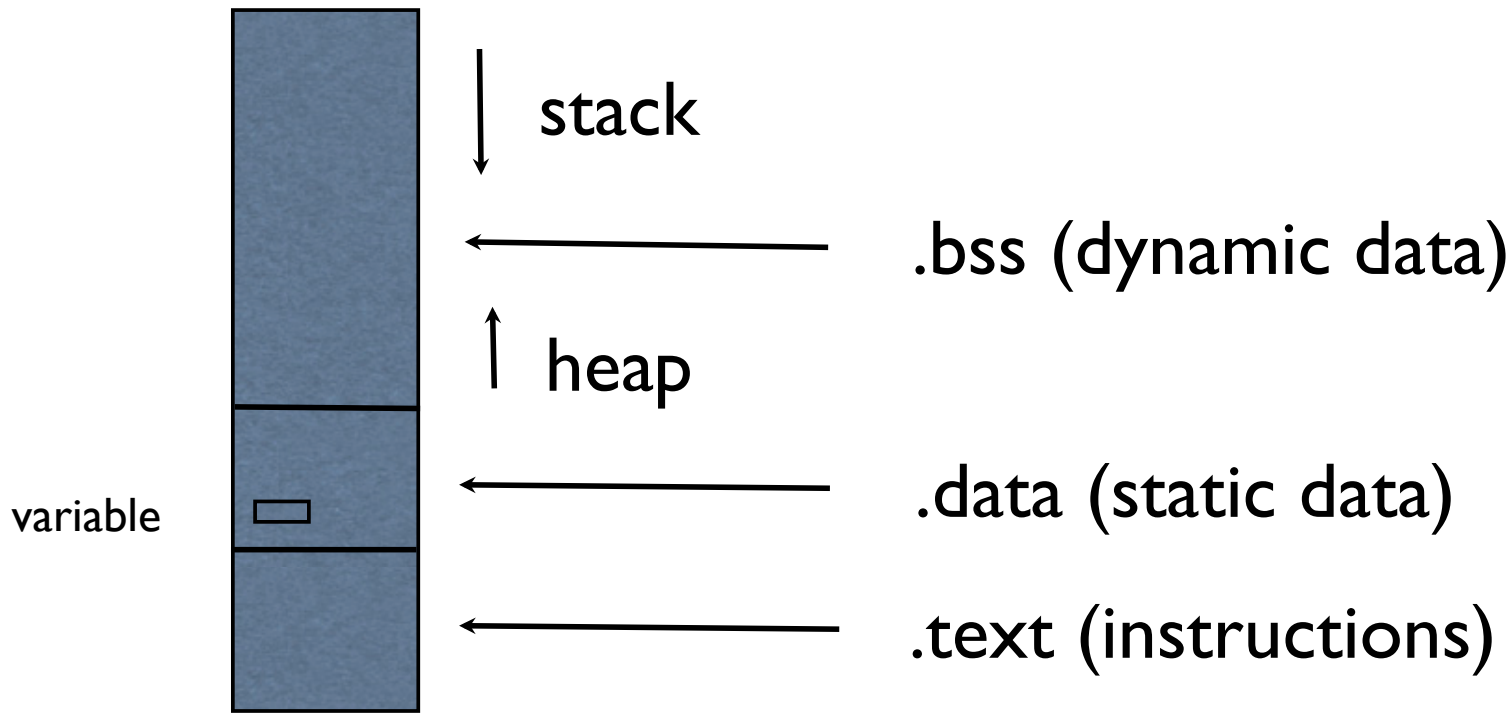
Brief History

- Agreed at Delft 2005
- Discussed at Fairfax 2006, at London 2007
- Las Vegas Compromise (core & more)
- Tokyo Further Compromise 2008
- Finally Agreed Las Vegas 2010
- Most recent draft of “more” is I2-I36r2

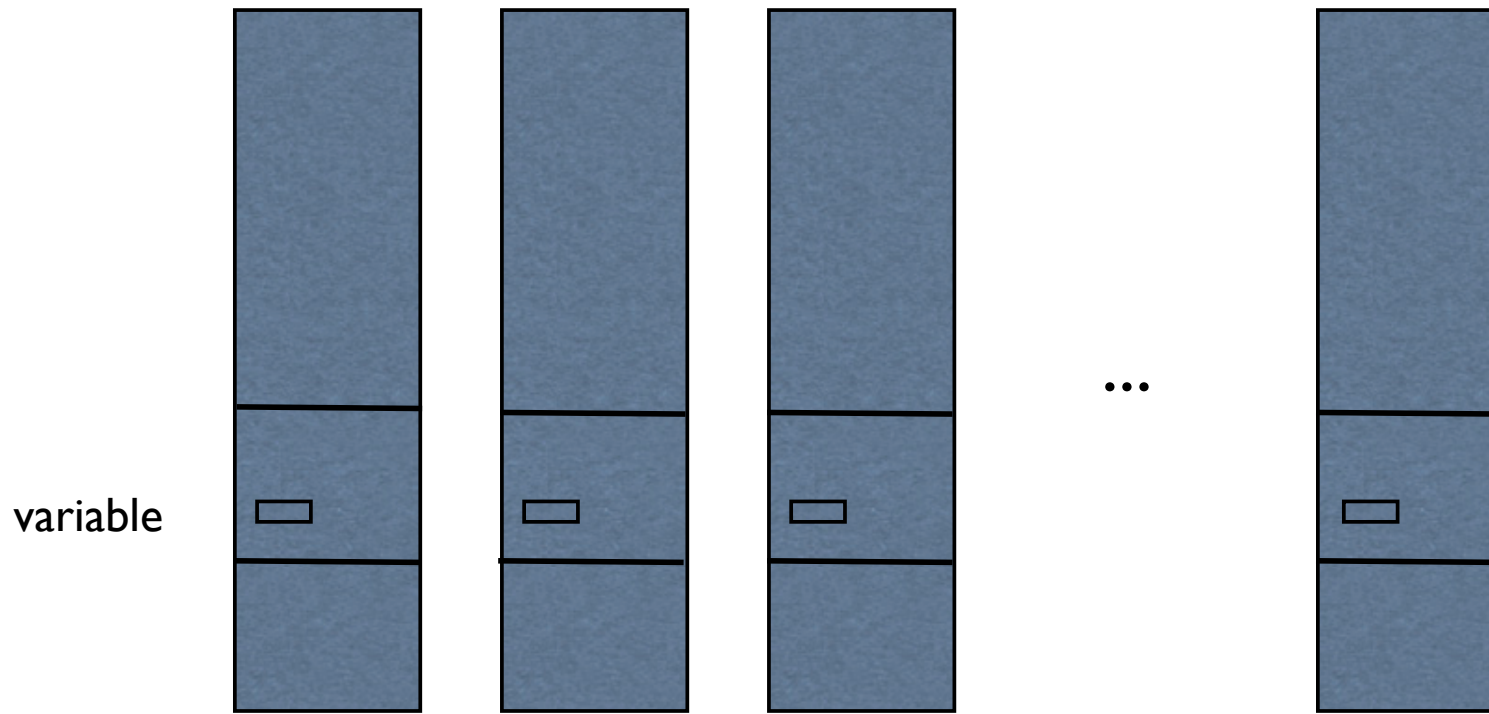
Coarray Concepts

- A program is treated as if it were replicated at the start of execution, each replication is called an image.
- Each image executes asynchronously.
- A coarray is indicated by trailing [].
- A data object without trailing [] is local.
- Explicit synchronization statements are used to maintain program correctness.

Memory Basics (Single Image)

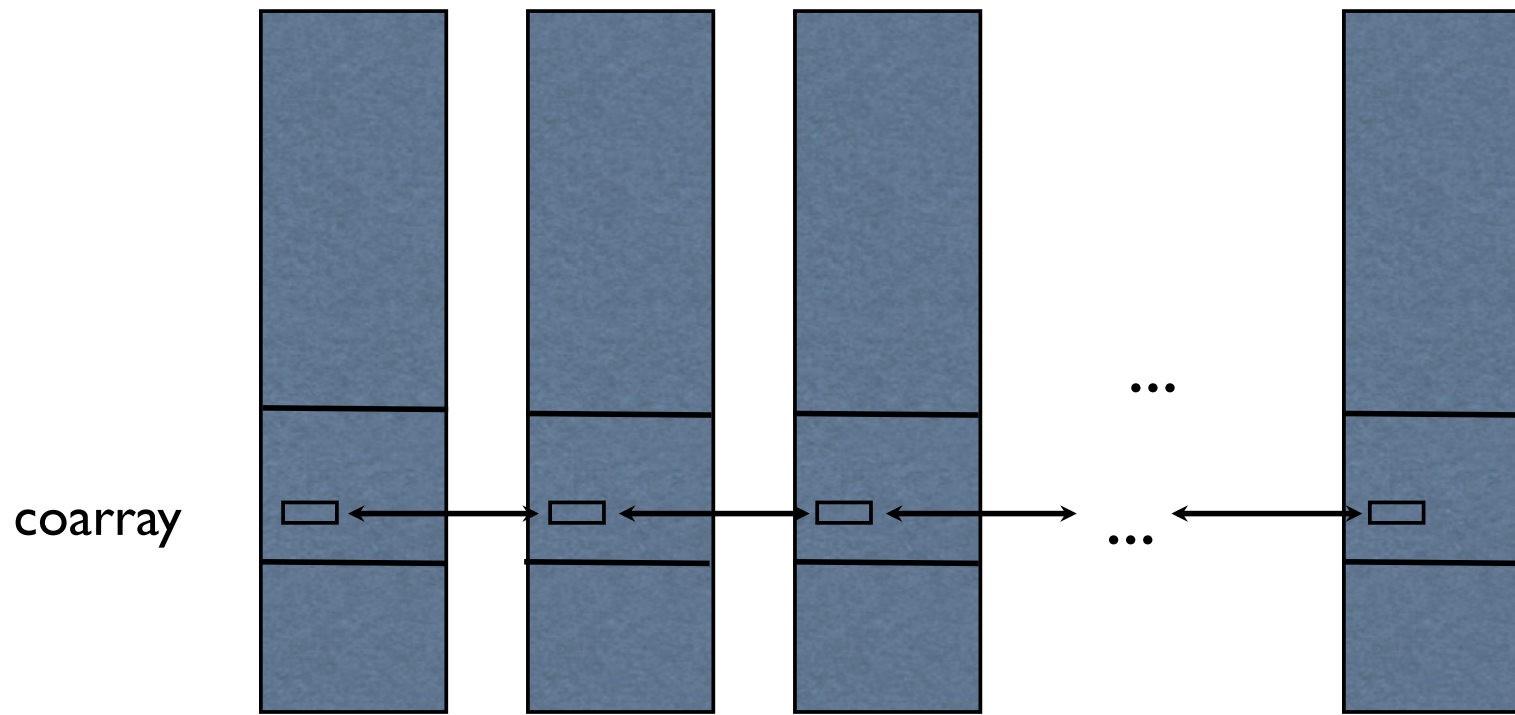


Memory Basics (Multiple Images)



Memory Basics

Coarrays



Core & More Compromise

- Reduce the coarray features in the Fortran 2008 standard (“core”)
- Add some of those features as a TS between revisions of the Fortran standard
- The standards committees are now writing the work list for the TS (“more”)
- The current draft of the work list is 12-136r2

12-136r2

- Teams
- Collective procedures
- Atomics
- Events
- Limited I/O support

Excluded by I2-I36r2

- coarray pointers or coscalars
- reduction of OOP restrictions on derived types with coarray components
- asymmetric allocatable variables
- asynchronous copy procedure

Defining Teams

- There is an “all images” team at the start of execution
- An image is a member of only one team at a time (within a nested hierarchy)
- Teams are identified by team variables of type `team_type` defined in `iso_fortran_env`
- Further teams are formed by a “split team” or a “form team” statement or intrinsic

Using Teams

- Supply a “with team ... end with team” or a “select team ... end select team” construct
- Within the construct, image numbers are relative to the specified team
- Collective activities are team-wide
- At the end of the team construct, the images revert to the previous team

Limitations on Teams

- Access to data outside the team is not allowed (how to synchronize?)
- Synchronization is team-wide
- Collective procedures apply team-wide
- Allocatable coarrays allocated during the team execution must be deallocated before the end of the team execution

Collective Procedures

- The currently expected set is
 - `co_sum`
 - `co_min`
 - `co_max`
 - `co_broadcast`
 - `co_reduce`

Applicability

- If teams are accepted, collective procedures are synchronous within the current team
- If not, some other means to specify the images partaking in a collective operation will be defined
- For example, a list of image numbers

Atomic Operations

- Fortran 2008 defines `atomic_define` and `atomic_ref`
- Proposed further atomic operations are:
 - `atomic_cas` (compare and swap)
 - `atomic_add`, `_and`, `_or`, `_xor`, `_swap`
 - `atomic_ax` (masked xor)

Events

- The original notify/query pair morphed into an event variable feature
- Event variables (a standard-defined type?)
- Define the operations:
 - post
 - query
 - wait

Limited Parallel I/O

- Opening a file for parallel I/O would
 - if teams are accepted, have a team= specifier on the open list
 - otherwise, have an images= specifier
 - synchronizes the team/images
 - be opened by all members of a team/images

More Parallel I/O

- Parallel I/O will be for direct access
- Explicitly closing the file is optional
 - if explicitly closed, all images must close the file and a synchronization occurs
- Use a flush statement to make data available to other images

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Thank you